

High dimensional geometry and Regularization

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This week

High dimensional Gaussians
Johnson Lindenstrauss Lemma

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High dimensional Gaussians
Johnson Lindenstrauss Lemma

Ridge and Lasso
Explanations
Compressive sensing
Matrix norms

High dimensional Gaussian

Multivariate Gaussian

$$f(\mathbf{x}) = \frac{1}{(2\pi)^{\frac{d}{2}} |\Sigma|^{\frac{1}{2}}} \exp\left(-(\mathbf{x} - \mu)^T \Sigma^{-1} (\mathbf{x} - \mu)\right)$$

$\mu = \mathbb{E}X$ (mean)

$\Sigma = \mathbb{E}(X - \mu)(X - \mu)^T$ (covariance)

High dimensional Gaussian

Multivariate Gaussian

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Where is the probability concentrated?

Concentration of measure

$$U \sim N(\mu, \sigma^2 I)$$

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Thin shell with width \sqrt{d}

For all $\delta > 0$,

$$\mathbb{P} \left(\|U - \mu\|^2 \leq \sigma^2 \left(d + 2\sqrt{d \ln \frac{1}{\delta}} \right) \right) \geq 1 - \delta$$

and

$$\mathbb{P} \left(\|U - \mu\|^2 \geq \sigma^2 \left(d - 2\sqrt{d \ln \frac{1}{\delta}} \right) \right) \geq 1 - \delta$$

Concentration of measure

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Around equator relative to any unit vector z

For all $\delta > 0$,

$$P\left(z^T(U - \mu) \leq \sigma \sqrt{2 \ln \frac{1}{\delta}}\right) \geq 1 - \delta$$

Johnson Lindenstrauss Lemma

Random projections preserve pairwise distances

For any ϵ and integer n , let $k = \frac{8 \ln n}{\epsilon^2}$. For all $z_1, \dots, z_n \in \mathbb{R}^d$, there exists $f : \mathbb{R}^d \rightarrow \mathbb{R}^k$ such that for all pairs z_i, z_j

$$\|f(z_i) - f(z_j)\|^2 \in (1 \pm \epsilon) \|z_i - z_j\|^2$$

These f can simply be random projections!

Applications of JL lemma

Regression in high dimensions

Some clustering problems
not always: GMM faster

Sketching and streaming algorithms

Learning mixtures of Gaussians

Cluster n points in \mathbb{R}^d into k clusters

Powerful and flexible model: Gaussian mixtures

$$X \sim \sum_{i=1}^k \pi_i \mathcal{N}(\mu_i, \Sigma_k)$$

Note: even common covariance $\Sigma_k = \Sigma$ versatile

Clustering in low dimensions, few clusters

k -means

choose centers μ_1, \dots, μ_k at random

assign each example to nearest mean

update centers and repeat prior step till convergence

Soft version: Expectation Maximization

Fits most likely GMM iteratively

For Gaussians, soft version of k -means

In high dimensions

Recall: most probability in $\mathcal{N}(\mu, \sigma^2 I)$ close to $\sigma\sqrt{d}$.

$$\mathbb{P} \left(\|X - \mu\|^2 \geq \sigma^2 \left(d - 2\sqrt{d \ln \frac{1}{\delta}} \right) \right) \geq 1 - \delta$$

Probability of finding a point near μ is $\exp(-\mathcal{O}(d))$

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Need $\exp(\mathcal{O}(d))$ points to even have a point $\leq \frac{1}{2}\sigma\sqrt{d}$!

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Most plausible data sizes: “few scattered specks of dust in an enormous void” (Dasgupta '99)

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Low dim algorithms need exponential in d examples

Key idea: Project into few dimensions

Linear projections: the projections are Gaussian too!

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it is possible PCA collapses components of the mixture on top of each other (or nearly so)

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For clustering, try Johnson-Lindenstrauss:

$\frac{1}{\epsilon^2} \log n$ projections retain all pairwise distances
projected space still too large
exponential in $\frac{1}{\epsilon^2} \log n$ is $n^{\frac{1}{\epsilon^2}}$

Key idea: Project into few dimensions

- Don't worry about retaining all pairwise distances
- $\mathcal{O}(\log k)$ projections
- retain distances between means
- push points closer to mean in each cluster!

Key idea: Project into few dimensions

Don't worry about retaining all pairwise distances

$\mathcal{O}(\log k)$ projections

retain distances between means

push points closer to mean in each cluster!

Series of recent results on several common examples in the low-d space, how they recover parameters in high-d space

Distance between means

If $\|\mu_1 - \mu_2\| > \Omega(d^{1/4})$, should expect to separate out clusters
Note that in this regime, the spheres are not disjoint

Yet we should expect all points in one cluster to be closer to each other than points in other clusters

Why Gaussian mixtures

In principle, GMs can model any continuous distribution

Two particular examples (projects):

- Asset returns (see paper on discord)

- fMRI (see paper on discord)

Gaussian random matrices

If A is a $k \times n$ matrix, entries iid Gaussian
rows, cols independently chosen Gaussian multivariate
satisfy something called the Restricted isometry property
all small subset of columns approximately orthogonal

Key property used in Compressed Sensing
extends the Shannon-Nyquist theorem
used to shorten MRI acquisition on conventional equipment,
network tomography, radio astronomy and optical interferometry
(aperture synthesis)

Compressed Sensing

If x^* is a S -sparse signal in \mathbb{R}^n

$y = Ax^*$ (ie k linear measurements of x)

If k is very small, can we still find x^* ?

Compare with Shannon-Nyquist sampling

Convex relaxation

$y = Ax$ is underdetermined

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infinite solutions
which solution to choose?

Finding sparsest solution too hard
NP-hard

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Compressed sensing to the rescue

Convex relaxation

Recall $y = Ax^*$ where $x^* \in \mathbb{R}^n$ is S -sparse

A : $k \times n$ random Gaussian matrix

Solve $\hat{x} = \arg \min \|x\|_1$ such that $Ax = y$

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Solution will coincide with the sparsest x provided

A satisfies the restricted isometry property

$k > S \log n$

Another project idea

Ridge and Lasso

Already noted, brief review

Matrix Completion

This will be our segue into next topic: NLP
Also a chance to learn about singular values

Matrix completion

M is a matrix of user preferences
(Netflix challenge 480189×17770 , today much larger)
user i , movie j , rating in position M_{ij}
most observations unknown
observations in set Ω

Complete M using Ω

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Complete M using Ω
no reason this is this even possible!

Reminiscent of compressed sensing
inferring sparse signal with very few measurements
equivalent of sparsity?

“Sparsity” in matrices

Rank of a matrix, Singular Value Decomposition

Rank 1 matrices

General rank k matrices

Singular value decomposition

Outer product expression

Autoencoders

Singular value decomposition/Genres

$$M = U\Sigma V^T$$

Rank = number of non-zero singular values

Low rank r : few (r) singular values

$$M = \sigma_1 u_1 v_1^T + \cdots + \sigma_r u_r v_r^T$$

u_i : how much does each person like genre i ? v_i : how would a person who likes genre i like each movie?

Matrix completion

Works when Ω is chosen uniformly at random

The rank of M is low

Singular value are incoherent with the standard basis
projecting the standard basis to the subspace
of singular vectors

Matrix completion: nuclear norm

$$X = \arg \min \|X\|_s$$

such that $X_{ij} = M_{ij}$, $ij \in \Omega$ where $\|X\|_s$ is the nuclear norm